

FAST COEFFICIENT COMPUTATION FOR ALGEBRAIC POWER SERIES IN POSITIVE CHARACTERISTIC

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ABSTRACT. We revisit Christol’s theorem on algebraic power series in positive characteristic and propose yet another proof for it. This new proof combines several ingredients and advantages of existing proofs, which make it very well-suited for algorithmic purposes. We apply the construction used in the new proof to the design of a new efficient algorithm for computing the N th coefficient of a given algebraic power series over a perfect field of characteristic p . It has several nice features: it is more general, more natural and more efficient than previous algorithms. Not only the arithmetic complexity of the new algorithm is linear in $\log N$ and quasi-linear in p , but its dependency with respect to the degree of the input is much smaller than in the previously best algorithm. Moreover, when the ground field is finite, the new approach yields an even faster algorithm, whose bit complexity is linear in $\log N$ and quasi-linear in \sqrt{p} .

1. INTRODUCTION

Given a perfect field k of characteristic $p > 0$, we address the following question: how quickly can one compute the N th coefficient f_N of an algebraic power series

$$f(t) = \sum_{n \geq 0} f_n t^n \in k[[t]],$$

where N is assumed to be a large positive integer?

As such, the question is rather vague; both the data structure and the computation model have to be stated more precisely. The algebraic series f is specified in $k[[t]]$ as some root of a polynomial $E(t, y)$ in $k[t, y]$, of degree $d = \deg_y E \geq 1$ and of height $h = \deg_t E$. To do this specification unequivocally, we will make several assumptions. First, we assume that E is *separable*, that is E and its derivative $E_y = \partial E / \partial y$ are coprime in $k(t)[y]$. Equivalently, the resultant $R(t) \in k[t]$ of E and of E_y with respect to y is nonzero; let ρ denote its valuation. Clearly, $\rho \leq \deg R \leq (2d-1)h$. Second, we assume that E is *irreducible*¹ in $k(t)[y]$. Note that both assumptions are satisfied if E is assumed to be the minimal polynomial of f and that irreducibility implies separability as soon as we know that E has at least one root in $k[[t]]$. The polynomial E might have several roots in $k[[t]]$. In order to specify uniquely its root f , we further assume that we are given $f_0, \dots, f_{2\rho}$ in k such that

$$E(t, f_0 + f_1 t + \dots + f_{2\rho} t^{2\rho}) \equiv 0 \pmod{t^{2\rho+1}}.$$

In other words, the data structure used to represent f is the polynomial E together with the initial coefficients $f_0, \dots, f_{2\rho}$. (Actually $\rho+1$ coefficients are enough to

¹The first assumption is not always implied by the second one, as exemplified by $E = y^p - t \in \mathbb{F}_p[t, y]$, and in general by any irreducible polynomial E in $k[t, y^p]$.

ensure the unicity of f . However $2\rho+1$ coefficients are needed to ensure its existence; for this reason, we will always assume the coefficients of f are given up to 2ρ .)

Under these assumptions, the classical Newton iteration [13] allows the computation of the first N coefficients of f in quasi-linear complexity $\tilde{O}(N)$. Here, and in the whole article (with the notable exception of Section 4), the algorithmic cost is measured by counting the number of basic arithmetic operations ($+$, $-$, \times , \div) and applications of the Frobenius map ($x \mapsto x^p$) and of its inverse ($x \mapsto x^{1/p}$) in the ground field k . The soft-O notation $\tilde{O}(\cdot)$ indicates that polylogarithmic factors in the argument are omitted. Newton's iteration thus provides a quasi-optimal algorithm to compute f_0, \dots, f_N . A natural question is whether faster alternatives exist for computing the coefficient f_N alone.

With the exception of the rational case ($d = 1$), where the N th coefficient can be computed in complexity $O(\log N)$ by binary powering, the most efficient algorithm currently known to compute f_N in characteristic 0 has complexity $\tilde{O}(\sqrt{N})$ [8]. It relies on baby step / giant step techniques, combined with fast multipoint evaluation.

Surprisingly, in positive characteristic p , a radically different approach leads to a spectacular complexity drop to $O(\log N)$. However, the big-O term hides a (potentially dramatic) dependency in p . The good behavior of this estimate with respect to the index N results from two facts. First, if the index N is written in base p as $(N_{\ell-1} \dots N_1 N_0)_p$, then the coefficient f_N is given by the simple formula

$$(1) \quad f_N = [(S_{N_{\ell-1}} \cdots S_{N_1} S_{N_0} f)(0)]^{p^\ell}$$

where the S_r ($0 \leq r < p$) are the *section* operators defined by

$$(2) \quad S_r \sum_n g_n t^n = \sum_n g_{pn+r}^{1/p} t^n.$$

Note that for the finite field \mathbb{F}_p the exponents p^ℓ in (1) and $1/p$ in (2) are useless, since the Frobenius map $x \mapsto x^p$ is the identity map in this case.

Second, by Christol's theorem [5, 6, 12], the coefficients sequence of an algebraic power series f over a perfect field k of characteristic $p > 0$ is *p-automatic*: this means that f generates a finite-dimensional k -vector space under the action of the section operators. Consequently, with respect to a fixed k -basis of this vector space, one can express: f as a column vector C , the section operators S_r as square matrices A_r ($0 \leq r < p$), and the evaluation at 0 as a row vector R . Then Formula (1) becomes

$$(3) \quad f_N = [R A_{N_{\ell-1}} \cdots A_{N_1} A_{N_0} C]^{p^\ell}.$$

Since ℓ is about $\log N$, and since the size of the matrices A_r does not depend on N , formula (3) yields an algorithm of complexity $O(\log N)$. This observation (for any p -automatic sequence) is due to Allouche and Shallit [1, Cor. 4.5]. However, this last assertion hides the need to first find the linear representation $(R, (A_r)_{0 \leq r < p}, C)$. As shown in [2], translating naively the p -automaticity in terms of linear algebra yields matrices A_r whose size can be about $d^2 h p^{2d}$. Thus, their precomputation has a huge impact on the cost with respect to the prime number p .

In the particular case of a prime field $k = \mathbb{F}_p$, and under the assumption $E_y(0, f_0) \neq 0$, this was improved in [2] by building on an idea originally introduced by Christol in [5]: one can compute f_N in complexity $\tilde{O}((h+d)^5 h p) + O((h+d)^2 h^2 \log N)$. Currently, this is the best complexity result for this task.

Contributions. We further improve the complexity result from [2] down to $\tilde{O}(d^2hp + d^\omega h) + O(d^2h^2 \log N)$ (Theorem 3.4, Section 3.2). Here ω is the exponent of matrix multiplication. In the case where k is a finite field, we propose an even faster algorithm, with bit complexity linear in $\log N$ and quasi-linear in \sqrt{p} (Theorem 4.1, Section 4). It is obtained by blending the approach in Section 3.2 with ideas and techniques imported from the characteristic zero case [8]. All these successive algorithmic improvements are consequences of our main theoretical result (Theorem 2.2, Section 2.2), which can be thought of as an effective version of Christol's Theorem (an in particular proves it).

2. EFFECTIVE VERSION OF CHRISTOL'S THEOREM

We keep the notations of the introduction. Christol's theorem is stated as follows.

Theorem 2.1 (Christol's theorem). *Let $f(t)$ in $k[[t]]$ be a formal series algebraic on $k(t)$, where k is a perfect field with positive characteristic. Then there exists a finite-dimensional k -vector space containing $f(t)$ and stable by the section operators.*

The aim of this section is to state and to prove an effective version of Theorem 2.1, upon which our forthcoming algorithms will be built. Our approach follows the initial treatment by Christol [5], which is based on Furstenberg's theorem [11]. We nevertheless get a more precise and more general result. The main new idea we use is a globalization argument in order to compare the section operators at 0 and at $f(t)$. This argument is based on the interpretation *via* the Frobenius operator given in §2.1 below; it is closely related to the Cartier operator used in a beautiful geometric proof due to Deligne [9], Speyer [15] and further studied by Bridy [4]. Unfortunately, this geometric proof does not fit our need since, for the algorithmic application we have in mind, it requires the computation of a ring of integers, which is potentially very costly (of complexity more than polynomial in d and h).

2.1. Frobenius and sections. Recall that the ground field k is assumed to be a perfect field of prime characteristic p , for example a finite field \mathbb{F}_q , where $q = p^s$. Let $K = k(t)$ be the field of rational functions over k and let $L = K[y]/(E)$ be the algebraic extension, of degree d over K , naturally attached with the root $f \in k[[t]]$ of $E(t, y) \in k[t, y]$.

Since k is a perfect field, the Frobenius endomorphism $F : k \rightarrow k$ defined by $x \mapsto x^p$ is an automorphism of k . It extends to a ring homomorphism, still denoted F , from $L[t^{1/p}]$ to L which raises an element of $L[t^{1/p}] = L^{1/p}$ to the power p . This homomorphism is an isomorphism and its inverse writes

$$(4) \quad F^{-1} = \sum_{r=0}^{p-1} t^{r/p} S_r,$$

where each S_r , with $0 \leq r < p$, maps L into itself.

The use in (4) of the same notation as in Formula (2) is not a mere coincidence. The algebraic series f provides an embedding of L into the field of Laurent series $k((t))$, which is the evaluation of an element $P(y)$ of L at the point $y = f(t)$. We will call $\text{eval}_f : L \rightarrow k((t))$ the map which sends $P(y)$ to $P(f(t))$. The Frobenius operator extends from k to $k((t))$, and the same holds for the sections S_r ($0 \leq r < p$). These extensions are exactly those of Formula (2). Thus the S_r 's in Formula (4)

are *global section* operators, while in Formula (2) they are *local* (around 0) *section* operators. Global and local operators are compatible; they are related by:

$$(5) \quad \mathbf{F} \circ \text{eval}_f = \text{eval}_f \circ \mathbf{F}, \quad \mathbf{S}_r \circ \text{eval}_f = \text{eval}_f \circ \mathbf{S}_r.$$

As for rational functions, the Frobenius operator and the section operators induce respectively a ring isomorphism \mathbf{F} from $K[t^{1/p}]$ onto K and maps σ_r ($0 \leq r < p$) from K onto K such that $\mathbf{F}^{-1} = \sum_{r=0}^{p-1} t^{r/p} \sigma_r$. The operators \mathbf{F} and \mathbf{S}_r ($0 \leq r < p$) are not K -linear but only k -linear. More precisely, for any λ in $K[t^{1/p}]$, μ in K , and z in L

$$(6) \quad \mathbf{F}(\lambda z) = \mathbf{F}(\lambda) \mathbf{F}(z) \quad \text{and} \quad \mathbf{S}_r(\mu z) = \sum_{s=0}^{p-1} t^{\lfloor \frac{r+s}{p} \rfloor} \sigma_s(\mu) \mathbf{S}_{r-s}(z),$$

in other words both \mathbf{F} and \mathbf{F}^{-1} are actually semi-linear.

2.2. The key theorem. Let $k[t, y]_{<h, <d}$ be the set of polynomials in $k[t, y]$ with height less than h and degree less than d .

Theorem 2.2. *For $P \in k[t, y]_{<h, <d}$ and for $0 \leq r < p$, there exists a (unique) polynomial Q in $k[t, y]_{<h, <d}$ such that*

$$(7) \quad \mathbf{S}_r \left(\frac{P}{E_y} \right) \equiv \frac{Q}{E_y} \pmod{E}.$$

The rest of this subsection is devoted to the proof of Theorem 2.2. Although mainly algebraic, the proof is based on the rather analytic remark that any algebraic function in $k(t)[f]$ can be obtained as the residue at $T = f$ of some rational function in $k(t, T)$ (see lemma 2.3). This idea was already used in Furstenberg [11], whose work has been inspiring for us. The main new insight of our proof is the following: we replace several small branches around zero are by one unique branch around a moving point. In order to make the argument work, we shall need further to relate the behaviour of the section operators around 0 and around the aforementioned moving point. This is where the reinterpretation of the \mathbf{S}_r 's in terms of Frobenius operators will be useful.

We consider the ring $\mathcal{H} = k((t))[[T]]$ of power series over $k((t))$. Its fraction field is the field $\mathcal{K} = k((t))((T))$ of Laurent series over $k((t))$. There is an embedding $k((t))[y] \rightarrow \mathcal{H}$ taking a polynomial in y to its Taylor expansion around f . Formally, it is simply obtained by mapping the variable y to $f+T$. It extends to a field extension $k((t))(y) \rightarrow \mathcal{K}$. We will often write $P(t, f+T)$ for the image of $P(t, y) \in k((t))(y)$ in \mathcal{K} . The field \mathcal{K} is moreover endowed with a *residue map* $\text{res} : \mathcal{K} \rightarrow k((t))$, defined by $\text{res} \left(\sum_{i=v}^{\infty} a_i T^i \right) = a_{-1}$ (by convention, $a_{-1} = 0$ if $v > -1$). It is clearly $k((t))$ -linear.

Lemma 2.3. *For any polynomial $P \in k((t))[y]$, the following equality holds:*

$$\text{res} \left(\frac{P(t, f+T)}{E(t, f+T)} \right) = \frac{P(t, f)}{E_y(t, f)}.$$

Proof. Since f is a simple root of E , the series $E(t, f+T)$ has a simple zero at $T = 0$. This means that it can be written $E(t, f+T) = T \cdot q(T)$ with $q \in \mathcal{H}$, $q(0) \neq 0$. Taking the logarithmic derivative with respect to T gives

$$\frac{E_y(t, f+T)}{E(t, f+T)} = \frac{1}{T} + \frac{q'(T)}{q(T)},$$

akin to [11, Formula (15), p. 276], from which we derive

$$\frac{P(t, f+T)}{E(t, f+T)} = \frac{g(T)}{T} + g(T) \frac{g'(T)}{q(T)},$$

where $g(T) = P(t, f+T)/E_y(t, f+T)$. Since $E_y(t, f+T)$ does not vanish at $T = 0$, the series $g(T)$ has no pole at 0. Therefore, the residue of $g(T)/T$ is nothing but $g(0)$. Besides the residue of the second summand $g(T)g'(T)/q(T)$ vanishes. All in all, the residue of $P(t, f+T)/E(t, f+T)$ is $g(0) + 0 = P(t, f)/E_y(t, f)$. \square

We now introduce analogues of section operators over \mathcal{K} . For this, we first observe that the Frobenius operator $x \mapsto x^p$ defines an isomorphism $F : \mathcal{K}[t^{1/p}, T^{1/p}] \rightarrow \mathcal{K}$. Moreover $\mathcal{K}[t^{1/p}, T^{1/p}]$ is a field extension of \mathcal{K} of degree p^2 . A basis of $\mathcal{K}[t^{1/p}, T^{1/p}]$ over \mathcal{K} is of course $(t^{r/p} T^{s/p})_{0 \leq r, s < p}$, but it will be more convenient for our purposes to use a different one. It is given by the next lemma.

Lemma 2.4. *The family $(t^{r/p} (f+T)^{s/p})_{0 \leq r, s < p}$ is a basis of $\mathcal{K}[t^{1/p}, T^{1/p}]$ over \mathcal{K} .*

Proof. For simplicity, we set $y = f+T \in \mathcal{K}$. We have:

$$(1 \quad y^{1/p} \quad \dots \quad y^{(p-1)/p}) = (1 \quad T^{1/p} \quad \dots \quad T^{(p-1)/p}) \cdot U$$

where U is the square matrix whose (i, j) entry (for $0 \leq i, j < p$) is $\binom{j}{i} f^{i/p}$. In particular, U is upper triangular and all its diagonal entries are equal to 1. Thus U is invertible and we are done. \square

For r and s in $\{0, 1, \dots, p-1\}$, we define the section operators $S_{r,s} : \mathcal{K} \rightarrow \mathcal{K}$ by

$$F^{-1} = \sum_{r=0}^{p-1} \sum_{s=0}^{p-1} t^{r/p} (f+T)^{s/p} S_{r,s}.$$

Clearly $S_{r,0}$ extends the operator $S_r : k((t)) \rightarrow k((t))$ defined by Eq. (2) and $S_{r,s}(g_1^p g_2) = g_1 S_{r,s}(g_2)$ for all $g_1, g_2 \in \mathcal{K}$. We observe moreover that the $S_{r,s}$'s stabilize the subrings $k((t))[y]$ and $k[t, y]$, since y corresponds to $f+T$.

Proposition 2.5. *The following commutation relation holds over \mathcal{K} :*

$$S_r \circ \text{res} = \text{res} \circ S_{r,p-1}.$$

Proof. Let us write $g \in \mathcal{K}$ as $g = \sum_{i=v}^{\infty} a_i T^i$ with $v \in \mathbb{Z}$ and $a_i \in k((t))$ for all $i \geq v$. Its image under F^{-1} can be expressed in two different ways as follows

$$F^{-1}(g) = \sum_{i=v}^{\infty} F^{-1}(a_i) T^{i/p} = \sum_{r=0}^{p-1} \sum_{s=0}^{p-1} t^{r/p} (f+T)^{s/p} S_{r,s}(g).$$

We identify the coefficient in $T^{-1/p}$. For doing so, we observe that the terms obtained with $s < p-1$ do not contribute, while the contribution of the term $t^{r/p} (f+T)^{(p-1)/p} S_{r,p-1}(g)$ is the residue of $t^{r/p} S_{r,p-1}(g)$. We then get

$$F^{-1}(a_{-1}) = \sum_{r=0}^{p-1} \text{res} \circ S_{r,p-1}(g) \cdot t^{r/p}.$$

Going back to the definition of S_r , we derive $S_r(a_{-1}) = \text{res} \circ S_{r,p-1}(g)$, from which the lemma follows. \square

Proof of Theorem 2.2. Let $P \in k[t, y]$ and $0 \leq r < p$. We set $Q = \mathcal{S}_{r, p-1}(PE^{p-1}) \in k[t, y]$. Combining Lemma 2.3 and Proposition 2.5, we derive the following equalities:

$$\begin{aligned} \mathcal{S}_r \left(\frac{P(t, f)}{E_y(t, f)} \right) &= \mathcal{S}_r \circ \text{res} \left(\frac{P(t, f+T)}{E(t, f+T)} \right) \\ &= \text{res} \circ \mathcal{S}_{r, p-1} \left(\frac{P(t, f+T)}{E(t, f+T)} \right) = \text{res} \left(\frac{Q(t, f+T)}{E(t, f+T)} \right) = \frac{Q(t, f)}{E_y(t, f)}. \end{aligned}$$

The stability of $k[t, y]/E(t, y)$ under \mathcal{S}_r follows using the fact that E is the minimal polynomial of f over $K = k(t)$. If we know in addition that P lies in $k[t, y]_{<h, <d}$ then PE^{p-1} is in $k[t, y]_{<ph, \leq p(d-1)}$ and therefore Q falls in $k[t, y]_{<h, <d}$ as well. Theorem 2.2 is proved. \square

Remark 2.6. It is possible to slightly vary the bounds on the degree and the height, and to derive in this way other stability statements. For example, starting from a polynomial $P(t, y)$ with $\deg_t P \leq h$ and $\deg_y P \leq d$, we have:

$$\mathcal{S}_r \frac{P(t, f)}{E_y(t, f)} = \frac{Q(t, f)}{E_y(t, f)}$$

with $\deg_t Q \leq h$ and $\deg_y P < d$. Moreover $\deg_t Q < h$ provided that $r > 0$.

Another remark in this direction is the following: if P has degree at most $d-2$, the section $\mathcal{S}_{r, p-1}(PE^{p-1})$ has degree no more than $d-2$ for any $r \in \{0, 1, \dots, p-1\}$. Indeed, PE^{p-1} has degree no more than $pd - 2 < p(d-1) + p - 1$. In other words, the subspace $k[t, y]_{<h, \leq d-2}$ is stable by the section operators \mathcal{S}_r ($0 \leq r < p$). Their matrices are then block triangular.

3. APPLICATION TO ALGORITHMICS

Theorem 2.2 exhibits an easy-to-handle finite dimension vector space which is stable under the section operators. In this section, we derive from it two efficient algorithms that compute the N th term of f in linear time with respect to $\log N$.

3.1. First algorithm: modular computations. The first algorithm we will design follows rather straightforwardly from Theorem 2.2. It consists of the following main steps:

- (1) we compute the matrix giving the action of the Frobenius F with respect to the “modified monomial basis” $\mathcal{B} = (y^j/E_y)_{0 \leq j \leq d-1}$;
- (2) we deduce the matrix of F^{-1} with respect to \mathcal{B} ;
- (3) we extract from it the matrices of the section operators \mathcal{S}_r ;
- (4) we compute the N th coefficient of f using Formula (1).

Let us be a bit more precise (though we will not give full details because we will design in §3.2 below an even faster algorithm). Let M be the matrix of F in the basis \mathcal{B} ; its j th column contains the coordinates of the vector $F(\frac{y^j}{E_y}) = \frac{y^{pj}}{E_y^p}$ in the basis \mathcal{B} , which are also the coordinates of y^{pj}/E_y^{p-1} in the monomial basis $(1, y, \dots, y^{d-1})$. It is easily seen that the matrix of F^{-1} with respect to \mathcal{B} is $F^{-1}(M^{-1})$, which is the matrix obtained by applying F^{-1} to each entry of M^{-1} .

We now discuss the complexity of the computation of M^{-1} . Thanks to Theorem 2.2, we know that its entries are polynomials of degree at most $h(p-1)$. However, this bound is not valid for the entries of M . Indeed, in full generality, the latter are rational fractions whose numerators and denominators have degrees of magnitude

Rearranging the terms, we finally find:

$$\begin{aligned} S_0(y) &= E_y^{-1} \cdot (4t^4 + (2t + 4t^2)y + (1 + t + 2t^2)y^2) \\ S_1(y) &= E_y^{-1} \cdot (4t^3 + (1 + 4t^3)y + (t + 4t^3)y^2) \\ S_2(y) &= E_y^{-1} \cdot ((2t^2 + 4t^3) + 3t^2y + (2 + 4t)y^2) \\ S_3(y) &= E_y^{-1} \cdot (4t + (t + 3t^2)y + (3 + 4t + 2t^2)y^2) \\ S_4(y) &= E_y^{-1} \cdot (1 + (3 + 3t)y + (4 + 3t)y^2). \end{aligned}$$

To conclude this example, suppose that we want to compute the 70th coefficient of f . Applying Eq. (1), we find that it is equal to the constant coefficient of $S_2 S_4 S_0 f$. Therefore we have to compute $S_2 S_4 S_0 y$. Repeating two times what we have done before, we end up with

$$S_2 S_4 S_0 y = E_y^{-1} \cdot ((2 + t^2) + (4 + 3t + 3t^3)y + (2 + 4t^2 + 2t^3)y^2).$$

Plugging $y = f$ in the above equality, we get $S_2 S_4 S_0 f = 2 + O(t)$. Consequently $f_{70} = 2$ and we are done.

Remark 3.2. In the above example, only the constant coefficient of f was needed to carry out the whole computation. This is related to the fact that $\text{Res}(E, E_y) = 4 + t + 3t^3 + t^4 + \dots$ has t -adic valuation 0. More generally if $\text{Res}(E, E_y)$ has t -adic valuation ρ , we will need the $\rho+1$ first coefficients of f since the final division by E_y will induce a loss of t -adic precision of ρ “digits”. This does not affect the complexity since $\rho \in O(dh)$.

3.2. Second algorithm: Hermite–Padé approximation. For obvious reasons related to the size of the computed objects, we cannot hope to achieve a complexity lower than linear with respect to p using the approach of §3.1. In order to improve the complexity of the whole computation, we return to Theorem 2.2. The key idea is to leap efficiently from the polynomial P to the polynomial Q in Formula (7).

Let $P = \sum_{i=0}^{d-1} a_i(t)y^i$ in $k[t, y]_{<h, <d}$ and $0 \leq r < p$. By Theorem 2.2, there exists $Q = \sum_{i=0}^{d-1} b_i(t)y^i$ in $k[t, y]_{<h, <d}$ such that $S_r(P/E_y) \equiv Q/E_y \pmod{E}$, or, equivalently,

$$(8) \quad S_r \left(\sum_{i=0}^{d-1} a_i(t) \frac{f(t)^i}{E_y(t, f(t))} \right) = \sum_{j=0}^{d-1} b_j(t) \frac{f(t)^j}{E_y(t, f(t))}.$$

The algorithmic question is to recover efficiently the b_i 's starting from the a_i 's. Identifying coefficients in Eq. (8) yields a linear system over k in the coefficients of the unknown polynomials b_i . This system has hd unknowns and an infinite number of linear equations. The point is that the following truncated version of Eq. (8)

$$(9) \quad S_r \left(\sum_{i=0}^{d-1} a_i(t) \frac{f(t)^i}{E_y(t, f(t))} \right) \equiv \sum_{j=0}^{d-1} b_j(t) \frac{f(t)^j}{E_y(t, f(t))} \pmod{t^{2dh}}$$

is sufficient to uniquely determine Q . This is a direct consequence of the following.

Lemma 3.3. *If Q in $k[t, y]_{<h, <d}$ satisfies $\frac{Q}{E_y}(t, f(t)) \equiv 0 \pmod{t^{2dh}}$, then $Q = 0$.*

Proof. The resultant $r(t)$ of $E(t, y)$ and $Q(t, y)$ with respect to y is a polynomial of degree at most $d(h-1) + h(d-1)$. On the other hand, we have a Bézout relation

$$E(t, y) u(t, y) + Q(t, y) v(t, y) = r(t),$$

Algorithm Nth coefficient via Hermite-Padé.

Input: A polynomial $E(t, y) = e_d(t)y^d + \cdots + e_0(t)$ and a truncation $g = f_0 + \cdots + O(t^{\rho+1})$ of a series f such that $E(t, g) = O(t^{\rho+1})$.

Output: The N th coefficient f_N of the series f .

1. Compute series expansions s_j of $f^j/E_y(t, f)$, $0 \leq j < d$, at precision $O(t^{2pdh})$.
 2. Expand $N = (N_{\ell-1} \dots N_0)_p$ with respect to the radix p .
 3. Set $g = y \in L$ written as $E_y^{-1} \cdot (-de_0 - (d-1)e_1y - \cdots - e_{d-1}y^{d-1})$.
 4. For $i = 0, 1, \dots, \ell - 1$,
 - (1) write $g = P(t, f)/E_y(t, f)$ as a linear combination of the s_j 's,
 - (2) compute the section $S_{N_i}(g)$ at precision $O(t^{2dh})$,
 - (3) recover Q such that $S_{N_i}(g) = Q/E_y$ by Hermite-Padé,
 - (4) put $g = Q/E_y$
 5. Substitute y by $f(t)$ in g and call $\bar{g}(t)$ the obtained result.
 6. Expand $\bar{g}(t)$ at precision $O(t)$.
 7. Set \bar{g}_0 to the constant coefficient of $\bar{g}(t)$.
 8. Return \bar{g}_0^p .
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where $u(t, y)$ and $v(t, y)$ are bivariate polynomials in $k[t, y]$. By evaluating the previous equality at $y = f(t)$ it follows that

$$r(t) \equiv Q(t, f(t)) v(t, f(t)) \equiv 0 \pmod{t^{2dh}}$$

holds in $k((t))$, and therefore $r = 0$. Thus E and Q have a non-trivial common factor; since E is irreducible, it must divide Q . But $\deg_y Q < \deg_y E$, so $Q = 0$. \square

Solving Eq. (9) amounts to solving a *Hermite-Padé approximation* problem. In terms of linear algebra, it translates into solving a linear system over k in the coefficients of the unknown polynomials b_i . This system has dh unknowns and $N = 2dh$ linear equations. Moreover, it has a very special shape: it has a quasi-Toeplitz structure, with displacement rank $\Delta = O(d)$. Therefore, it can be solved using fast algorithms for structured matrices [14, 3] in $\tilde{O}(\Delta^{\omega-1}N) = \tilde{O}(d^{\omega h})$ operations in k . These algorithms first compute a (quasi)-inverse of the matrix encoding the homogenous part of the system, using a compact data-structure called displacement generators (or, ΣLU representation); then, they apply it to the vector encoding the inhomogeneous part. The first step has complexity $\tilde{O}(\Delta^{\omega-1}N) = \tilde{O}(d^{\omega h})$, the second step has complexity $\tilde{O}(\Delta N) = \tilde{O}(d^2 h)$.

In our setting, we will need to solve $\log N$ systems of this type, each corresponding to the current digit of N in base p . An important feature is that these systems share the same homogeneous part, which only depends on the coefficients of the power series $s_j(t) = \frac{f^j}{E_y(t, f(t))}$ occurring on the right-hand side of (9). Only the inhomogeneous parts vary: they depend on the linear combination $\sum_{i=0}^{d-1} a_i(t)s_i(t)$. Putting these facts together yields Algorithm 1 and the following complexity result.

Theorem 3.4. *Let k be a perfect field with characteristic p . Let $E(t, y)$ be an irreducible polynomial in $k[t, y]$ of height h and degree d . Set $R(t)$ the resultant*

$\text{Res}_y(E, E_y)$. We assume that $R(t)$ is nonzero and we denote by ρ its t -adic valuation. We assume that we are given a polynomial $\bar{f}(t)$ such that $E(t, \bar{f}(t)) \equiv 0 \pmod{t^{2\rho+1}}$.

There there exists a unique series $f(t)$ congruent to $\bar{f}(t)$ modulo $t^{\rho+1}$ for which $E(t, f(t)) = 0$. Moreover, Algorithm 1 computes the N th coefficient of f for a cost of $\tilde{O}(d^2hp + d^\omega h) + O(d^2h^2 \log N)$ operations in k .

Proof. The first assertion is Hensel Lemma [10, Th. 7.3].

The precomputation of $s_j(t) = \frac{f(t)^j}{E_y(t, f(t))}$ modulo t^{2dhp} for $0 \leq j < d$ can be performed using Newton iteration, for a total cost of $\tilde{O}(d^2hp)$ operations in k . As explained above, this is enough to set up the homogeneous part of the quasi-Toeplitz systems; its inversion has cost $\tilde{O}(d^\omega h)$.

Let us turn to the main body of the computation, which depends on the index N . For each p -digit $r = N_i$ of N , we first construct the inhomogeneous part of the system. For this, we extract the coefficients of t^{pj+r} in $\sum_{i=0}^{d-1} a_i(t)s_i(t)$, for $0 \leq j < d$, for a total cost of $O(d^2h^2)$ operations in k . We then apply the inverse of the system to it, for a cost of $O(d^2h^2)$ (using a naive matrix vector multiplication). This is done $\ell \approx \log N$ times. The other steps of the algorithm have negligible cost. \square

4. IMPROVING THE COMPLEXITY WITH RESPECT TO p

As shown in Theorem 3.4, Algorithm 1 has a nice complexity with respect to the parameters d , h and $\log N$: it is polynomial with small exponents. However, the complexity with respect to p is not that good, as it is exponential in $\log p$, which is the relevant parameter. Thus, when p is large (say $> 10^5$), Algorithm 1 runs slowly and is no longer usable.

For this reason, it is important to improve the complexity with respect to p . In this section, we introduce some ideas to achieve this. More precisely, our aim is to design an algorithm whose complexity with respect to p and N is $\tilde{O}(\sqrt{p}) \cdot \log N$, and stays polynomial in all other relevant parameters. In the current state of knowledge, it seems difficult to lower further the exponent on p ; indeed, the question addressed in this paper is related to other intensively studied questions (e.g., counting points *via* p -adic cohomologies) for which the barrier $\tilde{O}(\sqrt{p})$ has not been overcome yet.

Notations and assumptions. We keep the notations of previous sections. We make one additional hypothesis: *the ground field k is a finite field*. We assume that k is presented as $(\mathbb{Z}/p\mathbb{Z})[X]/\pi(X)$ where π is an irreducible monic polynomial over $\mathbb{Z}/p\mathbb{Z}$ of degree s . We choose a monic polynomial $\hat{\pi} \in \mathbb{Z}[X]$ of degree s lifting π . We set $W = \mathbb{Z}_p[X]/\hat{\pi}(X)$ where \mathbb{Z}_p is the ring of p -adic integers.

The algorithm we are going to design is not algebraic in the sense that it does not only perform algebraic operations in the ground field k , but will sometimes work over W (or, more exactly, over finite quotients of W). For this reason, throughout this section, we will use bit complexity instead of algebraic complexity.

We use the notation $\text{poly}(n)$ for referring to a quantity whose growth is at most polynomial in n . The precise result we will prove states as follows.

Theorem 4.1. *Under the assumptions of Theorem 3.4 and the above extra assumptions, there exists an algorithm of bit complexity $\text{poly}(dhs)\tilde{O}(\sqrt{p}) \log N$ that computes the N th coefficient of f .*

If p is bounded by a (fixed) polynomial in d and h , then Theorem 4.1 has been proved already. In the sequel, we will then always assume that $p \gg d, h$.

Overview of the strategy. We reuse the structure of Algorithm 1 but speed up the computation of the $S_{N_i}(g)$'s. Precisely, in Algorithm 1, the drawback was the computation of the $\frac{f^j}{E_y(t,f)}$'s or, almost equivalently, the computation of $g = \frac{P(t,f)}{E_y(t,y)}$ at sufficient precision. However, only a few coefficients of g are needed since we are only interested in one of its sections. A classical method for avoiding this overhead is to find a (small) recurrence on the coefficients on $g = \sum_{n=0}^{\infty} g_i t^i$ of the form:

$$(10) \quad b_r(i)g_{i+r} + b_{r-1}(i)g_{i+r-1} + \cdots + b_1(i)g_{i+1} + b_0(i)g_i = 0.$$

We then unroll it using matrix factorials (for which fast algorithms are available in the literature [8]). Unrolling the recurrence is straightforward as soon as the leading coefficient $b_r(i)$ does not vanish. On the contrary, when $b_r(i) = 0$, the value of g_{i+r} cannot be deduced from the previous ones. Unfortunately, it turns out that $b_r(i)$ does sometimes vanish in our setting.

We tackle this issue by lifting everything over W and performing all computations over this ring. Divisions by p then become possible but induce losses of precision. We then need to control the p -adic valuation of the denominators, that are the p -adic valuations of the $b_r(i)$'s. We cannot expect to have a good control on them in full generality; even worse, we can build examples where $b_r(i)$ vanishes in W for some i . There exists nevertheless a good situation —the so-called *ordinary case*— where we can say a lot on the $b_r(i)$'s. With this extra input, we are able to lead our strategy to its end.

The general case reduces to the ordinary one using a change of origin, *i.e.* replacing t by $u+\alpha$ for some $\alpha \in k$. This change of origin does not seem to be harmless *a priori*. Indeed the Taylor expansion of g around α (the one we shall compute) has in general nothing to do with the Taylor expansion of g around 0 (the one we are interested in). The sections are nevertheless closely related (see Proposition 4.3). This “miracle” is quite similar to that we have already observed in Proposition 2.5 and again can be thought of as an avatar of the Cartier operator.

4.1. From algebraic equations to recurrences. We consider a bivariate polynomial $P(t, y) \in k[t, y]$ with $\deg_t P < h$ and $\deg_y P < d$. We fix moreover an integer r is the range $[0, p-1]$. Our aim is to compute $S_r\left(\frac{P(t,f)}{E_y(t,f)}\right)$ at precision $O(t^{2dh})$. Set $g = \frac{P(t,f)}{E_y(t,f)}$ and write $g = \sum_{i=0}^{\infty} g_i t^i$. By definition $S_r(g) = \sum_{j=0}^{\infty} g_{r+pj} t^j$, so that we have to compute the coefficients g_{r+pj} for $j < 2dh$.

We let L be the leading coefficient of $E(t, y)$ and R be the resultant of E and E_y . To begin with, we make the following assumption:

(H1): Both L and R have t -adic valuation 0.

As explained above, we now lift the situation over W . We choose a polynomial $\hat{E} \in W[t, f]$ of bidegree (d, h) lifting E . We define $\hat{E}_t = \frac{\partial \hat{E}}{\partial t}$, $\hat{E}_y = \frac{\partial \hat{E}}{\partial y}$. The assumption **(H1)** implies that the series f lifts uniquely to a series $\hat{f} \in W[[t]]$ such that $\hat{E}(t, \hat{f}) = 0$. We define \hat{L} as the leading coefficient of $\hat{E}(t, y)$ and set $\hat{R} = \text{Res}(\hat{E}, \hat{E}_y)$. We introduce the ring $W_K = W[t, (\hat{L}\hat{R})^{-1}]$. By **(H1)**, W_K embeds canonically into $W[[t]]$. We pick a polynomial $\hat{P} \in W[t, f]$ lifting P such that $\deg_t \hat{P} < h$ and $\deg_y \hat{P} < d$. We set $\hat{g} = \hat{P}(t, \hat{f})$.

We now compute a linear differential equation satisfied by \hat{g} . For this, we observe that the derivation $\frac{\partial}{\partial t} : W[[t]] \rightarrow W[[t]]$ stabilizes the subring $W_L = W_K[\hat{f}]$. Indeed

from the relation $\hat{E}(t, \hat{f}) = 0$, we get $\frac{\partial \hat{f}}{\partial t} = -\frac{\hat{E}_t(t, \hat{f})}{\hat{E}_y(t, \hat{f})}$. Thus $\frac{\partial \hat{f}}{\partial t} \in W_L$ because $\hat{E}_y(t, \hat{f})$ is invertible in W_L thanks to **(H1)**. Using additivity and the Leibniz relation, we finally deduce that $\frac{\partial}{\partial t}$ takes W_L to itself. In particular, all the successive derivatives of \hat{g} lie in W_L . On the other hand, we notice that W_L is free of rank d over W_K with basis $(1, \hat{f}, \dots, \hat{f}^{d-1})$. Let M be the $d \times d$ matrix whose j th column (for $0 \leq j < d$) contains the coordinates of $\frac{\partial^j \hat{g}}{\partial t^j}$ with respect to the above basis. Similarly let C be the column vector whose entries are the coordinates of $\frac{\partial^d \hat{g}}{\partial t^d}$. Let $\Delta_d = \det M$. We solve the system $MX = C$ using Cramer's formulae and find this way a linear differential equation of the form:

$$\Delta_d \frac{\partial^d \hat{g}}{\partial t^d} + \Delta_{d-1} \frac{\partial^{d-1} \hat{g}}{\partial t^{d-1}} + \dots + \Delta_1 \frac{\partial \hat{g}}{\partial t} + \Delta_0 \hat{g} = 0,$$

where the other Δ_i 's are defined as determinants as well. In particular, they all lie in W_K . Multiplying by the appropriate power of $\hat{L}\hat{R}$, we end up with a differential equation of the form:

$$(11) \quad \hat{a}_d \frac{\partial^d \hat{g}}{\partial t^d} + \hat{a}_{d-1} \frac{\partial^{d-1} \hat{g}}{\partial t^{d-1}} + \dots + \hat{a}_1 \frac{\partial \hat{g}}{\partial t} + \hat{a}_0 \hat{g} = 0$$

where the \hat{a}_i 's are polynomials in t . We can even be more precise. Indeed, following the above constructions, we find that all entries of M and C are rational functions whose degrees (of numerators and denominators) stay within $\text{poly}(dh)$. We then deduce that the degrees of the \hat{a}_i 's are in $\text{poly}(dh)$ as well. Furthermore, they can be computed for a cost of $\text{poly}(dh)$ operations in k , that is $\text{poly}(dh)\tilde{O}(s \log p)$ bit operations (recall that s denotes the degree of k over \mathbb{F}_p)

We write $\hat{g} = \sum_{i=0}^{\infty} \tilde{g}_i \frac{t^i}{i!}$. The differential equation (11) translates to a recurrence relation on the \tilde{g}_i 's of the shape:

$$(12) \quad \forall n \geq r, \quad \tilde{b}_0(n)\tilde{g}_n + \tilde{b}_1(n)\tilde{g}_{n-1} + \tilde{b}_2(n)\tilde{g}_{n-2} + \dots + \tilde{b}_r(n)\tilde{g}_{n-r} = 0$$

where the \tilde{b}_i 's are polynomials in n over W and r is at most $d + \max_i \deg \hat{a}_i$. In particular, $r \in \text{poly}(dh)$. Moreover it is easy to write down explicitly \tilde{b}_0 : it is the constant polynomial with value $\hat{a}_d(0)$.

4.2. The ordinary case. In order to take advantage of Eq. (12), we make the following extra assumption, corresponding to the so-called *ordinary case*:

(H2): $\hat{a}_d(0)$ does not vanish modulo p .

Under **(H2)**, $\tilde{b}_0(n) = \hat{a}_d(0)$ is invertible in W and there is no obstruction to unrolling the recurrence (12). Let us be more precise. We recall that we want to compute the values of g_{r+pj} for j up to $2dh$. Clearly g_n is the reduction modulo p of $\frac{\hat{g}_n}{n!}$. In order to get g_{r+pj} , we need to compute \tilde{g}_{r+pj} modulo p^{v+1} where v is the p -adic valuation of $(r + 2dhp)!$. Under our assumption that p is large enough compared to d and h , we get $v = 2dh$. We will then work over the finite ring $W' = W/p^{2dh+1}W$.

We first compute the r first coefficients of \hat{f} modulo p^{2dh+1} by solving the equation $\hat{E}(t, \hat{f}) = 0$ (using a Newton iteration for example). Since $r \in \text{poly}(dh)$, this computation can be achieved for a cost of $\text{poly}(dh)$ operations in W' , that is

$\text{poly}(dh)\tilde{O}(s \log p)$ bit operations. We then build the companion matrix:

$$M(n) = \begin{pmatrix} & & & 1 \\ & & \ddots & \\ & & & & 1 \\ \frac{-\tilde{b}_r(n)}{\hat{a}_d(0)} & \frac{-\tilde{b}_{r-1}(n)}{\hat{a}_d(0)} & \dots & \frac{-\tilde{b}_1(n)}{\hat{a}_d(0)} \end{pmatrix} \in M_r(W'[n]).$$

Obviously,

$$(\tilde{g}_{n-r+1} \ \tilde{g}_{n-r+2} \ \dots \ \tilde{g}_n)^T = M(n) \cdot M(n-1) \cdots M(r) \cdot (\tilde{g}_0 \ \tilde{g}_1 \ \dots \ \tilde{g}_{r-1})^T,$$

and computing \tilde{g}_n reduces to evaluating the matrix factorial $M(n) \cdot M(n-1) \cdots M(r)$. Using [8], the latter can be computed within $\text{poly}(dh)\tilde{O}(\sqrt{i})$ operations in W' , that is $\text{poly}(dh)\tilde{O}(\sqrt{i} \cdot s \log p)$ bit operations. All in all, we find that the g_{r+pj} 's ($0 \leq j < 2dh$) can all be computed for a cost of $\text{poly}(dh)\tilde{O}(s\sqrt{p})$ bit operations.

Plugging this input in Algorithm 1, we end up with an algorithm of bit complexity $\text{poly}(dh)\tilde{O}(s\sqrt{p}) \log N$. Theorem 4.1 is thus proved under the extra assumptions **(H1)** and **(H2)**.

4.3. Reduction to the ordinary case. We finally explain how **(H1)** and **(H2)** can be relaxed. The rough idea is to translate the origin at some point where they hold simultaneously.

The case of complete vanishing. Before proceeding, we need to deal with the case where the whole polynomial \hat{a}_d vanishes modulo p . This case is actually very special as shown by the next lemma, whose proof relies on the fact that for a generic g , the minimal-order (homogeneous) linear differential equation over $k(t)$ satisfied by g has order exactly d [7].

Lemma 4.2. *For a generic $g \in L = k(t)[y]/E(t, y)$, the reduction of \hat{a}_d modulo p does not vanish.*

We say that an element $g \in L$ is *good* if the corresponding \hat{a}_d does not vanish modulo p . Lemma 4.2 ensures that goodness holds generically. It then holds with high probability since we have assumed that the ground field k has a large cardinality. Consequently, even if we were unlucky and g was not good, we can produce with high probability a writing $g = g_1 + g_2$ where g_1 and g_2 are both good. Since moreover the section S_r is additive, we can recover $S_r(g)$ as $S_r(g_1) + S_r(g_2)$.

For this reason, in what follows we will assume safely that g is good.

Change of origin. Let $\hat{\alpha} \in W$ be such that $\hat{L}(\hat{\alpha}) \not\equiv 0 \pmod{p}$, $\hat{R}(\hat{\alpha}) \not\equiv 0 \pmod{p}$, $\hat{a}_d(\hat{\alpha}) \not\equiv 0 \pmod{p}$. Such an element exists (since k is assumed to be large) and can be found for a cost of $\text{poly}(dh)$ operations in k (e.g., by enumerating its elements).

We denote by $\alpha \in k$ the reduction of $\hat{\alpha}$ modulo p and assume that $\alpha \neq 0$ (otherwise, we are in the ordinary case). We perform the change of variable $\tau_\alpha : t \mapsto u + \alpha$. Note that τ_α induces isomorphisms $k(t) \rightarrow k(u)$ and $k(t)[y]/E(t, y) \rightarrow k(u)[y]/E(u - \alpha, y)$. Furthermore, the polynomial $E(\alpha, y) = 0$ has d simple roots in an algebraic closure of k . Let $f_{\alpha,0}$ be one of them. By construction, $f_{\alpha,0}$ lies in a finite extension ℓ of k of degree at most d . Moreover, by Hensel Lemma, $f_{\alpha,0}$ lifts uniquely to a solution:

$$f_\alpha = f_{\alpha,0} + f_{\alpha,1}u + \dots + f_{\alpha,i}u^i + \dots \in \ell[[u]]$$

of the equation $E(u-\alpha, y) = 0$. We emphasize that the morphism $k(t)[y]/E(t, y) \rightarrow k(u)[y]/E(u-\alpha, y)$ does *not* extend to a mapping $k((t)) \rightarrow k((u))$ sending f to f_α . The next diagram summarizes the previous discussion:

$$\begin{array}{ccc}
\begin{array}{c} \text{S}_r \\ \curvearrowright \\ k((t)) \end{array} & & \begin{array}{c} \text{S}_{r,u} \\ \curvearrowleft \\ k((u)) \end{array} \\
\downarrow & & \downarrow \\
\frac{k(t)[y]}{E(t, y)} & \xrightarrow{\tau_\alpha} & \frac{k(u)[y]}{E(u-\alpha, y)} \\
\downarrow & & \downarrow \\
k(t) & \xrightarrow{\tau_\alpha} & k(u)
\end{array}$$

Here S_r and $\text{S}_{r,u}$ refer to the section operators defined in the usual way. We observe that they stabilize the subfields $\frac{k(t)[y]}{E(t, y)}$ and $\frac{k(u)[y]}{E(u-\alpha, y)}$ respectively since they can alternatively be defined by the relations:

$$(13) \quad \begin{array}{l} \text{over } \frac{k(t)[y]}{E(t, y)}: \quad \text{F}^{-1} = \sum_{r=0}^{p-1} t^{r/p} \text{S}_r \\ \text{over } \frac{k(u)[y]}{E(u-\alpha, y)}: \quad \text{F}^{-1} = \sum_{r=0}^{p-1} u^{r/p} \text{S}_{r,u} \end{array}$$

where F is the Frobenius map (see also Eq. (4)).

Proposition 4.3. *The commutation $\text{S}_{p-1,u} \circ \tau_\alpha = \tau_\alpha \circ \text{S}_{p-1}$ holds over $\frac{k(t)[y]}{E(t, y)}$.*

Proof. Clearly τ_α commutes with the Frobenius because it is a ring homomorphism. From the relations (13), we then derive $\sum_{r=0}^{p-1} u^{r/p} \text{S}_{r,u} \circ \tau_\alpha = \sum_{r=0}^{p-1} (u+\alpha)^{r/p} \tau_\alpha \circ \text{S}_r$. Identifying the coefficients in $u^{\frac{p-1}{p}}$, we get the announced result. \square

We emphasize that the other section operators $\text{S}_{r,*}$ (with $r < p-1$) *do not commute* with τ_α : the above phenomenon is specific to the index $p-1$. However, we can relate S_r and $\text{S}_{p-1,u}$ as follows.

Corollary 4.4. *For all $g \in \frac{k(t)[y]}{E(t, y)}$, we have $\text{S}_r(g) = \tau_\alpha^{-1} \circ \text{S}_{p-1,u} \circ \tau_\alpha(t^{p-1-r}g)$.*

Proof. This follows from Proposition 4.3 and from $\text{S}_r(g) = \text{S}_{p-1}(t^{p-1-r}g)$. \square

A modified recurrence. In order to use Corollary 4.4, we need to check that $\tau_\alpha(t^{p-1-r}g)$ fits the ordinary case. Recall the differential equation satisfied by \hat{g} ,

$$\hat{a}_d \frac{\partial^d \hat{g}}{\partial t^d} + \hat{a}_{d-1} \frac{\partial^{d-1} \hat{g}}{\partial t^{d-1}} + \cdots + \hat{a}_1 \frac{\partial \hat{g}}{\partial t} + \hat{a}_0 \hat{g} = 0.$$

We set $r' = p-1-r$ and $\hat{G} = t^{r'} \hat{g}$. Applying Leibniz formula to $\hat{g} = t^{-r'} \hat{G}$, we get:

$$\frac{\partial^j \hat{g}}{\partial t^j} = \sum_{i=0}^j (-1)^i \binom{j}{i} r'(r'+1) \cdots (r'+i-1) t^{-r'-i} \frac{\partial^j \hat{G}}{\partial t^{j-i}},$$

from which we derive the following differential equation satisfied by \hat{G} :

$$\sum_{0 \leq i \leq j \leq d} (-1)^i \hat{a}_j \binom{j}{i} r'(r'+1) \cdots (r'+i-1) t^{-r'-i} \frac{\partial^{j-i} \hat{G}}{\partial t^{j-i}}.$$

Reorganizing the terms and multiplying by $t^{r'+d}$, we end up with:

$$(14) \quad \sum_{j=0}^d \sum_{i=0}^{d-j} (-1)^i \hat{a}_{i+j} \binom{i+j}{i} r'(r'+1) \cdots (r'+i-1) t^{d-i} \frac{\partial^j \hat{G}}{\partial t^j}.$$

Set $W_{L,u} = W[u, y]/\hat{E}(u+\alpha, y)$ and define the ring homomorphism $\tau_{\hat{\alpha}} : W_L \rightarrow W_{L,u}$, $t \mapsto u+\hat{\alpha}$, $y \mapsto y$. Clearly $\tau_{\hat{\alpha}}$ lifts τ_{α} . Applying $\tau_{\hat{\alpha}}$ to Eq. (14) and noticing that $\frac{\partial}{\partial t} = \frac{\partial}{\partial u}$, we obtain:

$$\sum_{j=0}^d \sum_{i=0}^{d-j} (-1)^i \hat{\tau}_{\hat{\alpha}}(a_{i+j}) \binom{i+j}{i} r'(r'+1) \cdots (r'+i-1) (u+\hat{\alpha})^{d-i} \frac{\partial^j \tau_{\hat{\alpha}}(\hat{G})}{\partial u^j}.$$

Conclusion. The leading term of the latest differential equation (obtained only with $j = d$ and $i = 0$) is $\tau_{\hat{\alpha}}(\hat{a}_d) (u+\hat{\alpha})^d$. Its value at $u = 0$ is then $\hat{a}_d(\hat{\alpha}) \hat{\alpha}^d$, which is not congruent to 0 modulo p by assumption. Moreover the other coefficients are polynomials in u whose degrees stay within $\text{poly}(dh)$. Therefore, we can apply the techniques of §4.2 and compute $S_{p-1,u}(\tau_{\alpha}G)$ at precision $O(u^{2dh})$ for a cost of $\text{poly}(dh)\tilde{O}(s\sqrt{p})$ bit operations. As explained in §3.2, we can reconstruct $S_{p-1,u}(\tau_{\alpha}G)$ as an element of $k[u, y]/E(u+\alpha, y)$ for a cost of $\text{poly}(dh)$ operations in k using Hermite–Padé approximations. Thanks to Corollary 4.4, it now just remains to apply τ_{α}^{-1} to get $S_r(g)$. This last operation can be performed for a cost of $\text{poly}(dh)$ operations in k as well. All in all, we are able to compute $S_r(g)$ for a total bit complexity of $\text{poly}(dh)\tilde{O}(s\sqrt{p})$. Repeating this process $\log N$ times, we obtain the complexity announced in Theorem 4.1.

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